

# Safe Dynamic Memory Management in Ada and SPARK

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# Why Try To Verify Use of Pointers?

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- ▶ Automatic storage management
- ▶ Control “unknown” aliasing of names
- ▶ Use pointers in SPARK for formal verification

## How?

- ▶ Implement a variant of pointer *Ownership*
  - ↳ Inspired from Rust
  - ↳ Concurrent-Read-Exclusive-Write (CREW) policy

# Quick Reminder 1/2

- ▶ **Named** types:

```
type Int_Ptr is access Integer;  
X : Int_Ptr;
```

- ▶ **Anonymous** types:

```
Y : access Integer;
```

- ▶ **General** access types

```
type Int_Cst_Ptr is access constant Integer;  
type Int_Cst_Ptr is access all Integer;
```

- ▶ **Pool-specific** access types

↪ No `general_access_modifier` appears

## Quick Reminder 2/2

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- ▶ Access types

```
X : Int_Ptr;
```

- ▶ Composite types

```
type Rec is record  
  Data : Int_Ptr;  
end record
```

- ▶ By-copy types

Parameter passed by copy

- ▶ By-reference types

Parameter passed by reference: a view on the actual parameter

# Motivating Example: Swap Pointers

```
1 type Int_Ptr is access Integer;  
2  
3 procedure Swap(X_Param, Y_Param : in out Int_Ptr) is  
4   Tmp : Int_Ptr := X_Param;  
5   begin  
6     X_Param := Y_Param;  
7     Y_Param := Tmp;  
8   end Swap;  
9  
10 X : Int_Ptr := new Integer;  
11 Y : Int_Ptr := new Integer;  
12  
13 Swap(X, Y);
```

Dangling refs?

Storage leaks?

Correct result?

# Pointer Ownership: Overview

**Idea:** No more than one “owning” pointer to a given object

- Constraints:**
- ▶ Composite types are by-reference types
    - ↪ Always passed by reference
  - ▶ Access types are pool-specific types
    - ↪ Cannot point to stack

**Goal:** Automatic storage management

# Pointer Ownership: Overview

**Idea:** No more than one “owning” pointer to a given object

- Constraints:**
- ▶ Composite types are by-reference types
    - ↪ Always passed by reference
  - ▶ Access types are pool-specific types
    - ↪ Cannot point to stack

**Goal:** Automatic storage management

## Operations

- ▶ Move → complete transfer of the ownership
- ▶ Borrow → temporary transfer of the ownership
- ▶ Observe → no owning object

# Pointer Ownership: Overview

**Idea:** No more than one “owning” pointer to a given object

- Constraints:**
- ▶ Composite types are by-reference types
    - ↪ Always passed by reference
  - ▶ Access types are pool-specific types
    - ↪ Cannot point to stack

**Goal:** Automatic storage management

## Operations

- ▶ Move
- ▶ Borrow
- ▶ Observe

## Objects states

- ▶ Unrestricted  $\Rightarrow$  Read Write
- ▶ Observed  $\Rightarrow$  Read Only
- ▶ Borrowed  $\Rightarrow$  No



# Contents

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- 3 Observing Operations
- 4 Formal Verification in SPARK
- 5 Conclusion

# Moving Operations

# Moving Access Values

**When:** Assignment to **named**

- ▶ variables or return objects
- ▶ parameters of mode out/in-out

**Example:**

```
Y : Int_Ptr;  
X : Int_Ptr := Y;
```

## Conditions

- ▶ X, Y of named type
- ▶ X, Y unrestricted

## Results

- ▶ X unrestricted
- ▶ Old storage of X deallocated
- ▶ Y unrestricted, **null**

# Moving Composite Types

**When:** Assignment to

- ▶ variables or return objects
- ▶ ~~parameters of mode out/in-out~~

**Example:**

```
R : Rec := (...);  
S : Rec := (...);  
S := R;
```

## Conditions

- ▶ R, S unrestricted

## Results

- ▶ S unrestricted
- ▶ Old S components deallocated
- ▶ R unrestricted; components **null**

# Borrowing Operations

# Borrowing Access Values

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**When:** Initializing

- ▶ `in` parameters
- ▶ stand-alone **anonymous** objects
- ▶ constants

of an access-to-**variable** type

# Borrowing Access Values

## When: Initializing

- ▶ in parameters
- ▶ stand-alone **anonymous** objects
- ▶ constants

of an access-to-**variable** type

## Example:

```
procedure f(X_Param : in Int_Ptr);  
  f(X);
```

### Conditions

- ▶ X\_Param of mode in;  
 access-to-variable type
- ▶ X unrestricted

### Results

- ▶ X\_Param unrestricted
- ▶ X borrowed

# Borrowing Access Values

## When: Initializing

- ▶ `in` parameters
- ▶ **anonymous** stand-alone objects
- ▶ constants

of an access-to-**variable** type

## Example:

```
X : access Integer := Y
```

### Conditions

- ▶ X of an anonymous access-to-variable type
- ▶ Y unrestricted

### Results

- ▶ X unrestricted
- ▶ Y borrowed



# Borrowing Access Values

**When:** Initializing

- ▶ in parameters
- ▶ **anonymous** stand-alone objects
- ▶ constants

of an access-to-**variable** type

**Example:**

```
X : access Integer := Y
```

## Conditions

- ▶ X of an anonymous access-to-variable type
- ▶ Y unrestricted

## Results

- ▶ X unrestricted
- ▶ Y borrowed

# Borrowing Composite Objects

**When:** Passing parameters of mode out or in-out

**Example:**

```
procedure f(X_Param : in out Rec);  
  f(X);
```

## Conditions

- ▶ X\_Param of mode in-out
- ▶ X passed by-reference

## Results

- ▶ X\_Param unrestricted
- ▶ X borrowed

# Observing Operations

# Observing Access Values

**When:** Initializing

- ▶ `in` parameters
- ▶ **anonymous** stand-alone objects

of an access-to-**constant** type

**Example:**

```
X : access constant Integer := Y;
```

## Conditions

- ▶ X of an anonymous access-to-constant type
- ▶ Y unrestricted or observed

## Results

- ▶ X observed → Read Only
- ▶ Y observed → Read Only

# Observing Access Values

## When: Initializing

- ▶ `in` parameters
- ▶ **anonymous** stand-alone objects

of an access-to-**constant** type

## Example:

```
X : access constant Integer := Y;
```

### Conditions

- ▶ X of an anonymous access-to-constant type
- ▶ Y unrestricted or observed

### Results

- ▶ X observed → Read Only
- ▶ Y observed → Read Only

# Observing Composite Types

**When:** Initializing

- ▶ **constant** stand-alone objects
- ▶ parameters of mode `in`

**Example:**

```
procedure f(X_Param : in Rec);  
  f(X);
```

## Conditions

- ▶ `X_Param` of mode `in`
- ▶ `X` passed by-reference

## Results

- ▶ `X_Param` observed  $\rightarrow$  RO
- ▶ `X` observed  $\rightarrow$  RO

# Observing Composite Types

## When: Initializing

- ▶ **constant** stand-alone objects
- ▶ parameters of mode **in**

## Example:

```
procedure f(X_Param : in Rec);  
  f(X);
```

### Conditions

- ▶ X\_Param of mode **in**
- ▶ X passed by-reference

### Results

- ▶ X\_Param observed  $\rightarrow$  RO
- ▶ X observed  $\rightarrow$  RO

## Example Cont'd: Swap Pointers

```
1 type Int_Ptr is access Integer;
2
3 procedure Swap(X_Param, Y_Param : in out Int_Ptr) is
4   Tmp : Int_Ptr := X_Param ;
5 begin
6   X_Param := Y_Param ;
7   Y_Param := Tmp ;
8 end Swap;
9
10 X : Int_Ptr := new Integer;
11 Y : Int_Ptr := new Integer;
12
13 Swap( X , Y );
```

Tmp is the new owning object

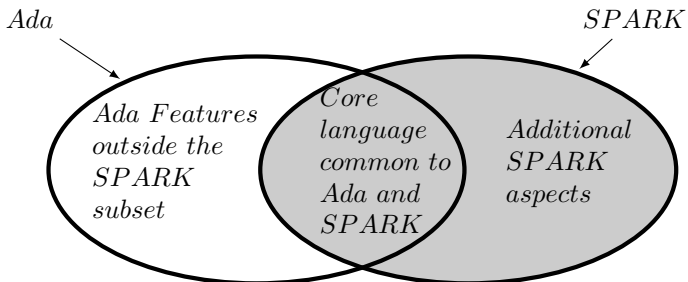
↪ No dangling reference, cannot dereference the old value of X\_Param



# Formal Verification in SPARK

# SPARK - What it is?

- ▶ A programming language
  - ▶ A subset of Ada, designed for static verification
  - ▶ Additional features to enhance program specification



- ▶ A set of program verification tools

# Why Aliasing Matters in SPARK?

```
1  type Int_Ptr is access Integer;
2
3  procedure Add_One(X_Param, Y_Param : in Int_Ptr) with
4     Post => X_Param.all = X_Param.all 'Old + 1
5     and Y_Param.all = Y_Param.all 'Old + 1
6  is
7  begin
8     X_Param.all := X_Param.all + 1;
9     Y_Param.all := Y_Param.all + 1;
10 end Add_One;
```

# Why Aliasing Matters in SPARK?

```

1  type Int_Ptr is access Integer;
2
3  procedure Add_One(X_Param, Y_Param : in Int_Ptr) with
4     Post => X_Param.all = X_Param.all'Old + 1
5     and Y_Param.all = Y_Param.all'Old + 1
6  is
7  begin
8     X_Param.all := X_Param.all + 1;
9     Y_Param.all := Y_Param.all + 1;
10 end Add_One;

```

If SPARK ignored aliasing:

```

1  X : Int_Ptr := new Integer'(3);
2  (...)
3
4  Add_One (X, X);
5  pragma Assert (X.all = 4);  -- incorrect assertion

```

# With Ownership Types: Alias

```

spark_proof.adb
1  procedure Spark_Proof with SPARK_Mode => On is
2
3  type Int_Ptr is access Integer;
4
5  procedure Add_One(X Param, Y Param : Int_Ptr) with
6    Post => X_Param.all = X_Param.all'Old + 1
7    and Y_Param.all = Y_Param.all'Old + 1
8  is
9  begin
10     X_Param.all := X_Param.all + 1;
11     Y_Param.all := Y_Param.all + 1;
12  end Add_One;
13
14     X : Int_Ptr := new Integer'(3);
15     Y : Int_Ptr := new Integer'(4);
16  begin
17
18     Add_One(X, X);
19     pragma Assert (Y.all = 5);
20
21  end Spark_Proof;

```

Spark\_Proof.Add\_One

Messages Locations Breakpoints /bin/bash

```

gnatprove -P/home/maalej/tests/spark/pointers/default.gpr -j0 --mode=flow --ide-progress-bar -u spark_proof.adb
Phase 1 of 2: generation of Global contracts ...
Phase 2 of 2: analysis of data and information flow ...
spark_proof.adb:18:15: insufficient permission for "X"
spark_proof.adb:18:15: expected state "unrestricted" at least, got "borrowed"
gnatprove: error during analysis of data and information flow
[2018-06-08 16:06:19] process exited with status 1, elapsed time: 00.81s

```

# With Ownership Types: Alias

---

```
spark_proof.adb:18:15: insufficient permission for "X"  
spark_proof.adb:18:15: expected state "unrestricted" at least, got "borrowed"
```

# With Ownership Types: Alias Free

```
spark_proof.adb
1 procedure Spark_Proof with SPARK_Mode => On is
2
3   type Int_Ptr is access Integer;
4
5   procedure Add_One(X_Param, Y_Param : Int_Ptr) with
6     Post => X_Param.all = X_Param.all'Old + 1
7     and Y_Param.all = Y_Param.all'Old + 1
8   is
9   begin
10    X_Param.all := X_Param.all + 1;
11    Y_Param.all := Y_Param.all + 1;
12  end Add_One;
13
14  X : Int_Ptr := new Integer'(3);
15  Y : Int_Ptr := new Integer'(4);
16 begin
17
18  Add_One(X, Y);
19  pragma Assert (Y.all = 5);
20
21 end Spark_Proof;
```

Spark\_Proof.Add\_One

Messages Locations Breakpoints

Q-Info

- Builder results (7 of 18 items in 1 file)
- spark\_proof.adb (7 of 18 items)
- 6:14 info: postcondition proved
- 6:44 info: overflow check proved
- 7:40 info: overflow check proved
- 10:12 info: initialization of "X" proved
- 10:15 info: initialization of "Y" proved
- 19:19 info: initialization of "Y" proved
- 19:19 info: assertion proved

# Conclusion



# Conclusion

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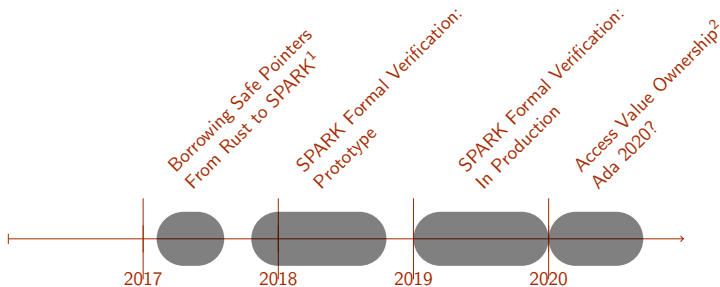
## Pointer Ownership Approach

- ▶ Inspired from Rust
- ▶ Safe pointers w.r.t to CREW policy: full ownership (read/write access); partial ownership (read-only access)

## Pointer Ownership Goals

- ▶ For Ada
  - ▶ No storage leaks
  - ▶ No dangling references
- ▶ For SPARK
  - ▶ No hidden aliasing  $\mapsto$  Can verify correctness of algorithms

# Supporting Pointers in SPARK: Steps



<sup>1</sup>Georges-Axel Jaloyan, Yannick Moy, and Andrei Paskevich. Borrowing Safe Pointers From Rust in SPARK. 2017. URL: <https://arxiv.org/abs/1805.05576>.

<sup>2</sup>AdaCore. Access value ownership and parameter aliasing. 2018. URL: <http://www.ada-auth.org/cgi-bin/cvsweb.cgi/ai12s/ai12-0240-1.txt>.

Questions?